# Portable Data Mining on Azure and HPC Platform

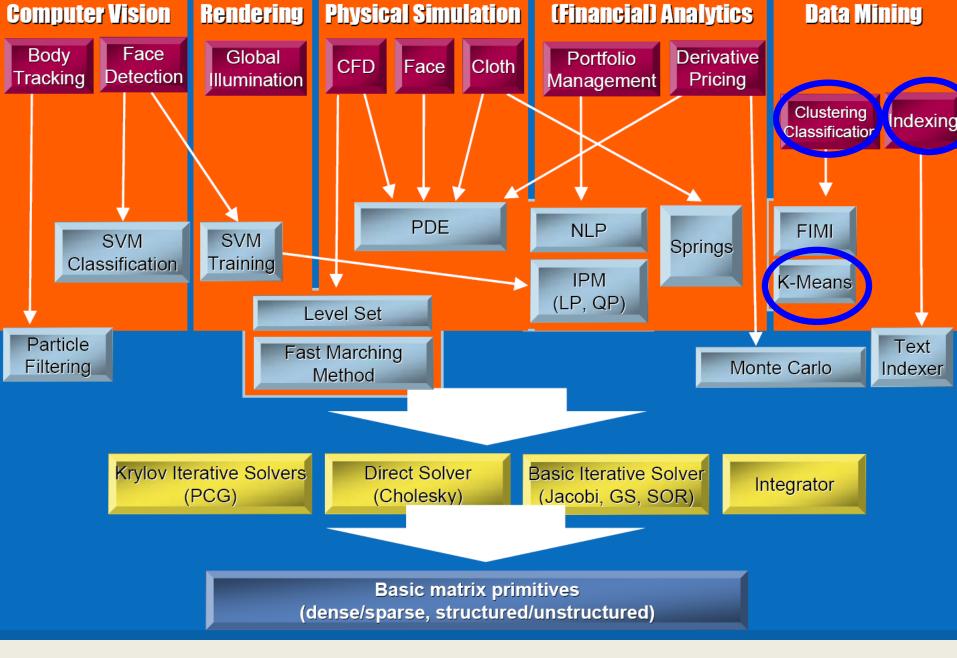
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### **Data Intensive Iterative Applications**

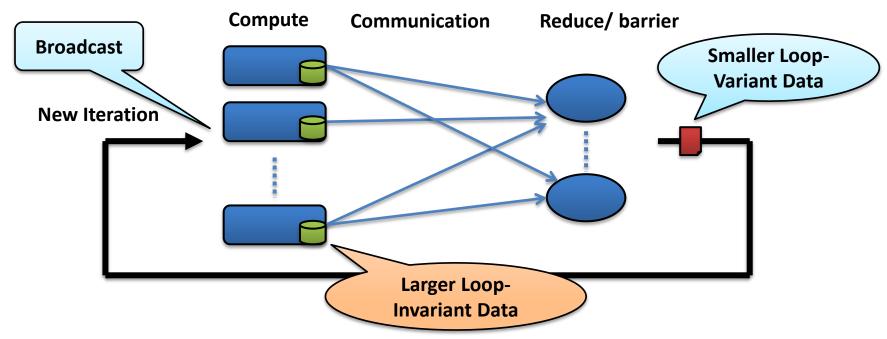
- Growing class of applications
  - Clustering, data mining, machine learning & dimension reduction applications Expectation Maximization
  - Driven by data deluge & emerging computation fields
  - Lots of scientific applications

```
\begin{split} & k \notin 0; \\ & \text{MAX} \leftarrow \text{maximum iterations} \\ & \delta^{[0]} \leftarrow \text{initial delta value} \\ & \textbf{while (} k < \text{MAX_ITER || } f(\delta^{[k]}, \delta^{[k-1]}) ) \\ & \textbf{foreach datum in data} \\ & \beta[\text{datum}] \leftarrow \text{process (datum, } \delta^{[k]}) \\ & \textbf{end foreach} \\ & \delta^{[k+1]} \leftarrow \text{combine}(\beta[]) \\ & k \leftarrow k+1 \\ & \textbf{end while} \end{split}
```



Intel's Application Stack

#### **Data Intensive Iterative Applications**



- Common Characteristics
- Compute (map) followed by LARGE communication Collectives (reduce)

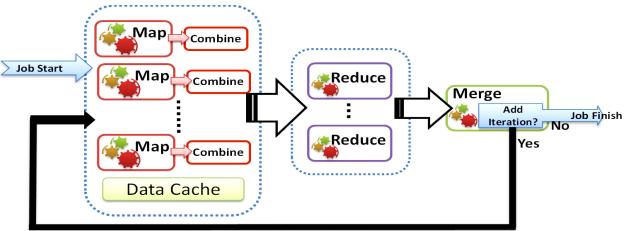
#### **Iterative MapReduce**

MapReduceMerge



 Extensions to support additional broadcast (+other) input data

> Map(<key>, <value>, list\_of <key,value>) Reduce(<key>, list\_of <value>, list\_of <key,value>) Merge(list\_of <key,list\_of<value>>,list\_of <key,value>)



Hybrid scheduling of the new iteration

# Parallel Data Analysis using Twister

#### Data mining and Data analysis Applications

- Next Generation Sequencing
- 🔮 Image processing
- 🔮 Search Engine
- ····

#### 🔮 Algorithms

- Multidimensional Scaling (MDS)
- 🔮 Clustering (K-means)
- 🔮 Indexing
- 🥘 ....

# Challenges

- Traditional MapReduce and classical parallel runtimes cannot solve iterative algorithms efficiently
  - Hadoop: Repeated data access to HDFS, no optimization to data caching and data transfers
  - MPI: no natural support of fault tolerance and programming interface is complicated

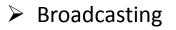
#### Interoperability

## **Current and Future Work**

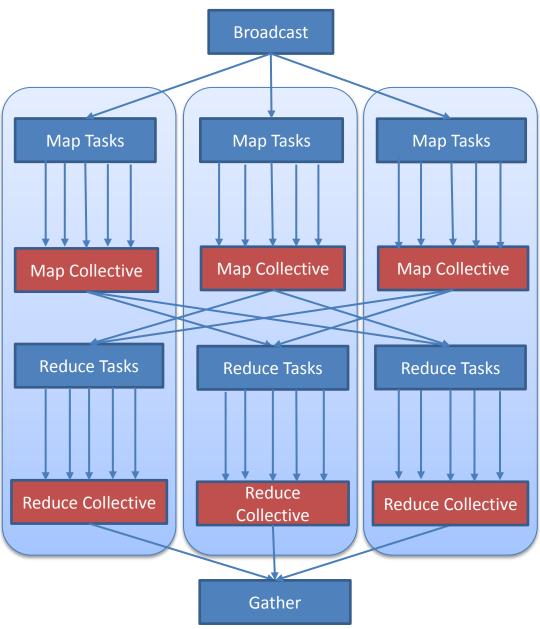
#### Collective Communication

- Fault tolerance
- Distributed Storage
- 🔮 High level language

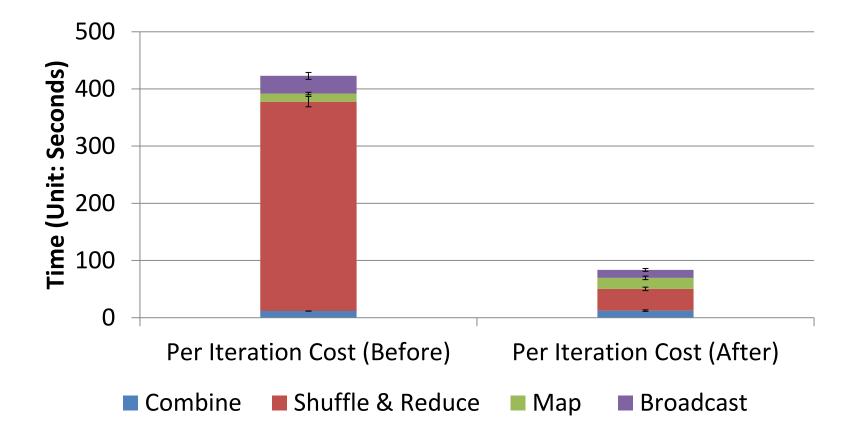
### **Twister Collective Communications**



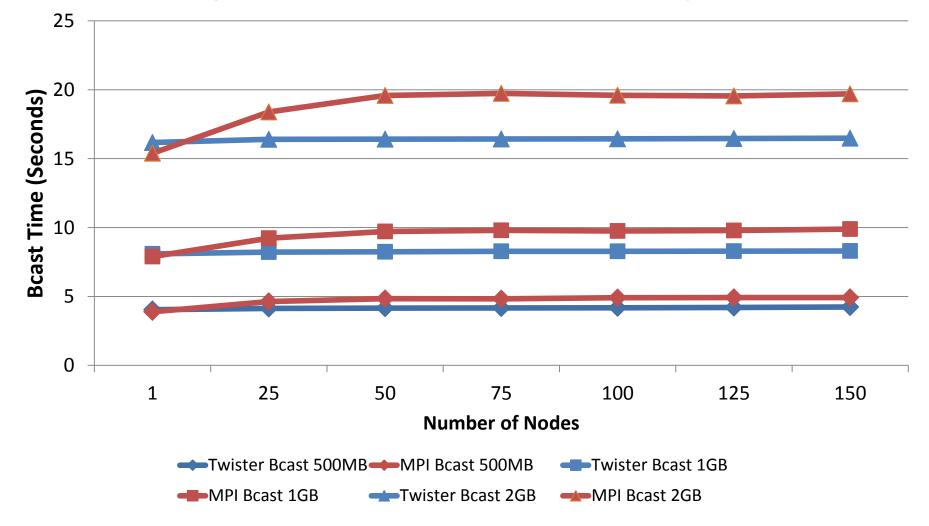
- Data could be large
- Chain & MST
- Map CollectivesLocal merge
- Reduce Collectives
  - Collect but no merge
- Combine
  - Direct download or Gather



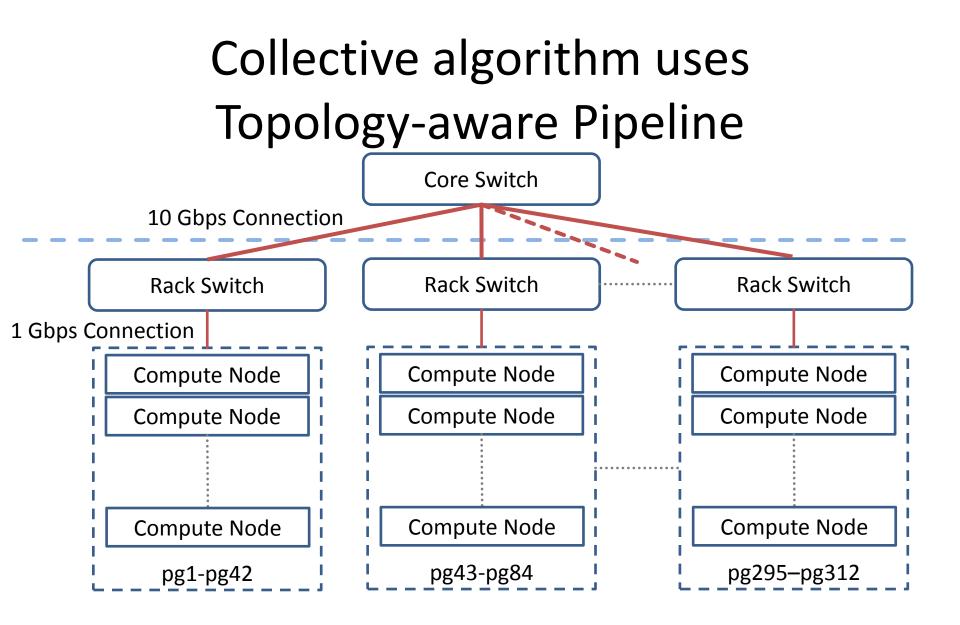
#### **Twister Broadcast Comparison** One-to-All vs. All-to-All implementations



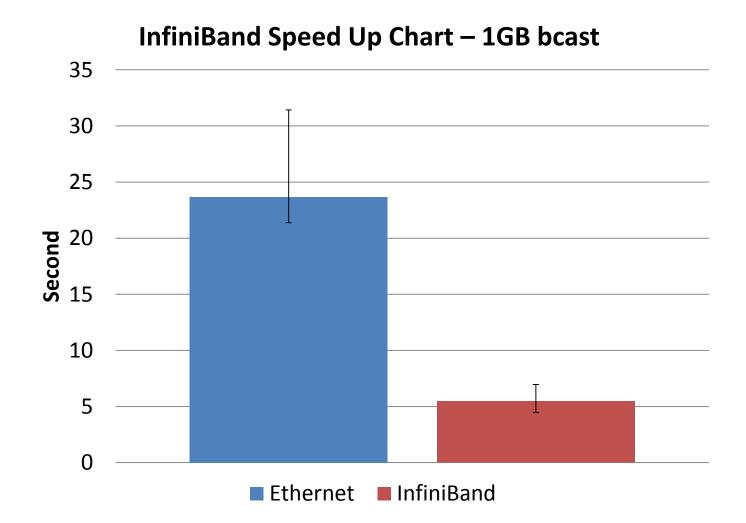
# Bcast Byte Array on PolarGrid (Fat-Tree Topology with 1Gbps Ethernet): Twister v. MPI (OpenMPI)



We are optimizing Collectives needed in data mining

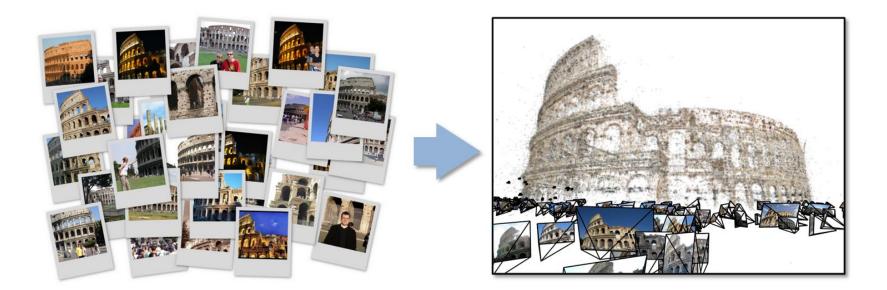


#### **Twister Broadcast Comparison:** Ethernet vs. InfiniBand (Oak Ridge)



## Data Intensive Kmeans Clustering

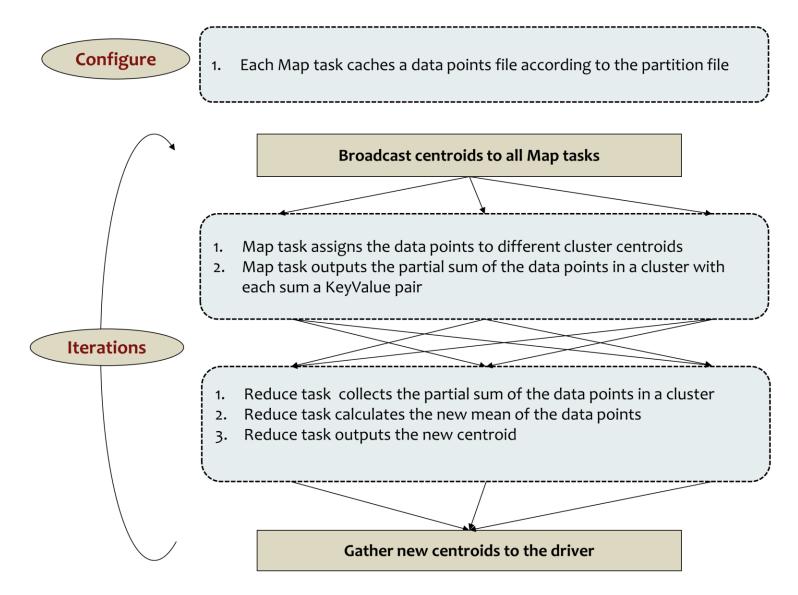
- *Image Classification:* **1.5 TB**; 500 features per image;10k clusters 1000 Map tasks; 1GB data transfer per Map task node



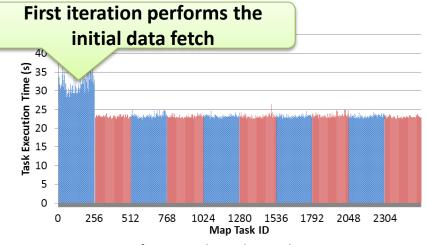
# High Dimensional Data

- K-means Clustering algorithm is used to cluster the images with similar features.
- In image clustering application, each image is characterized as a data point with 512 dimensions. Each value ranges from 0 to 255.
- Currently, we are able able to process 10 million images with 166 machines and cluster the vectors to 1 million clusters
  - Need 180 million images
- Improving algorithm (Elkan) and runtime (Twister Collectives)

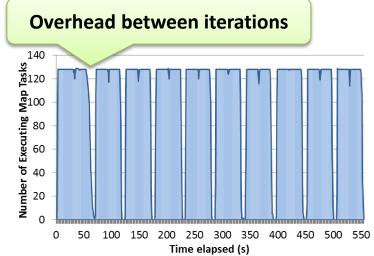
## **Twister Kmeans Clustering**



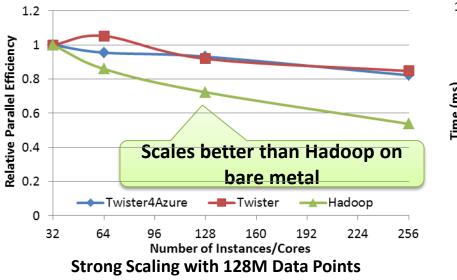
### Kmeans Clustering on Twister4Azure

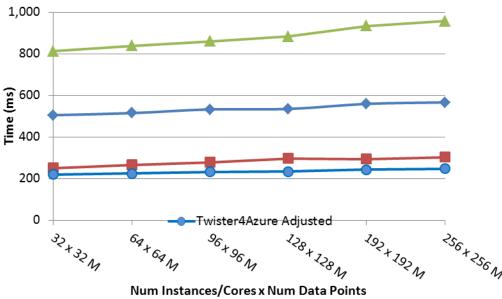


**Task Execution Time Histogram** 



#### Number of Executing Map Task Histogram





Weak Scaling

# **Triangle Inequality and Kmeans**

- Dominant part of Kmeans algorithm is finding nearest center to each point O(#Points \* #Clusters \* Vector Dimension)
- Simple algorithms finds min over centers c: d(x, c) = distance(point x, center c)
- But most of d(x, c) calculations are wasted as much larger than minimum value
- Elkan (2003) showed how to use triangle inequality to speed up using relations like

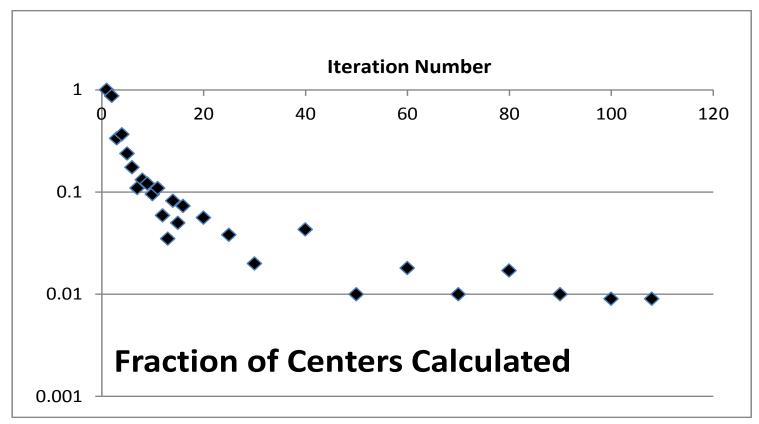
d(x, c2) >= d(x, c2-last) - d(c2, c2-last) and

d(x, c2) >= d(c1, c2) - d(x,c1)

c2-last position of center at last iteration; c1 c2 two centers

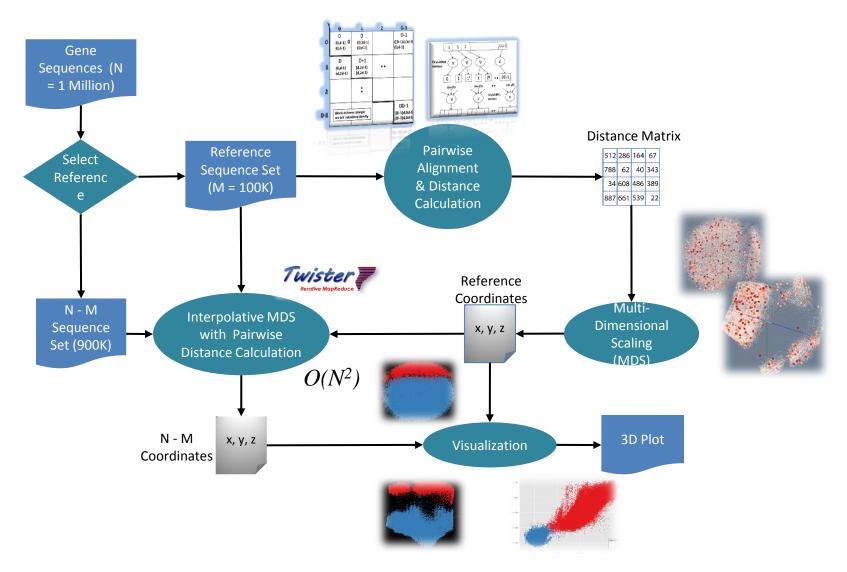
- So compare estimate of d(x, c2) with d(x, c1) where c1 is nearest cluster at last iteration
- Complexity reduced by a factor = Vector Dimension and so this important in clustering high dimension spaces such as social imagery with 500 or more features per image

# **Early Results on Elkan's Algorithm**



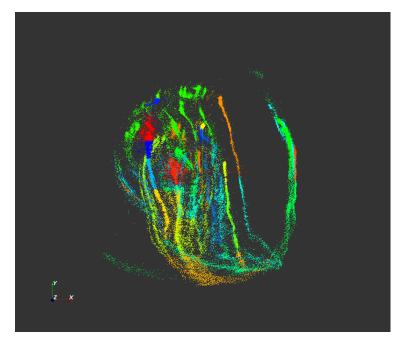
- Graph shows fraction of distances d(x, c) that need to be calculated each iteration for a test data set
- Only 5% on average of distance calculations needed
- 200K points, 124 centers, Vector Dimension 74

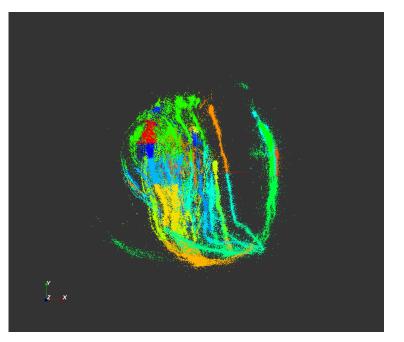
## **Bioinformatics Pipeline**



# **Million Sequence Challenge**

- Input DataSize: 680k
- Sample Data Size: 100k
- Out-Sample Data Size: 580k
- Test Environment: PolarGrid with 100 nodes, 800 workers.

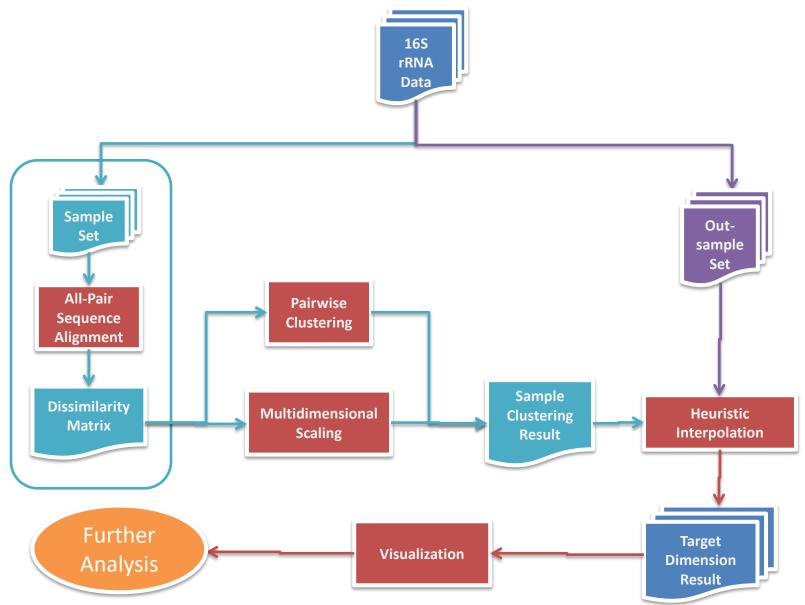




#### 100k sample data

680k data

#### DACIDR (A Deterministic Annealing Clustering and Interpolative Dimension Reduction Method) Flow Chart



## **Dimension Reduction Algorithms**

#### • Multidimensional Scaling (MDS) [1]

- Given the proximity information among points.
- Optimization problem to find mapping in target dimension of the given data based on pairwise proximity information while minimize the objective function.
- Objective functions: STRESS (1) or SSTRESS (2)

$$\sigma(\boldsymbol{X}) = \sum_{i < j \le N} w_{ij} (d_{ij}(\boldsymbol{X}) - \delta_{ij})^2$$
(1)

$$\sigma^{2}(\mathbf{X}) = \sum_{i < j \le N} w_{ij} [(d_{ij}(\mathbf{X}))^{2} - (\delta_{ij})^{2}]^{2}$$
(2)

- Only needs pairwise distances  $\delta_{ij}$  between original points (typically not Euclidean)
- d<sub>ij</sub>(X) is Euclidean distance between mapped (3D) points

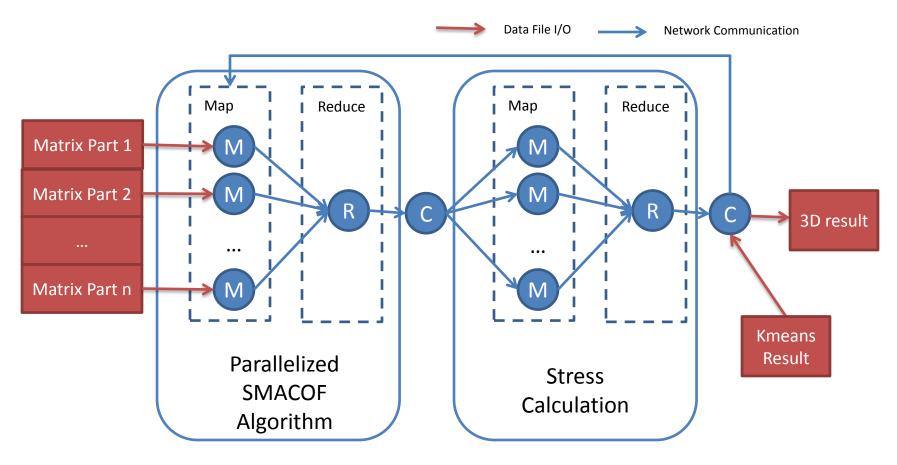
- Generative Topographic Mapping (GTM) [2]
  - Find optimal K-representations for the given data (in 3D), known as K-cluster problem (NP-hard)
  - Original algorithm use EM method for optimization
  - Deterministic Annealing algorithm can be used for finding a global solution
  - Objective functions is to maximize log-

$$\mathcal{L}(\boldsymbol{W},\beta) = \sum_{j=1}^{N} \ln \left\{ \frac{1}{K} \sum_{i=1}^{K} \mathcal{N}(\boldsymbol{x}_{j} | f(\boldsymbol{z}_{i}; \boldsymbol{W}), \beta) \right\}$$

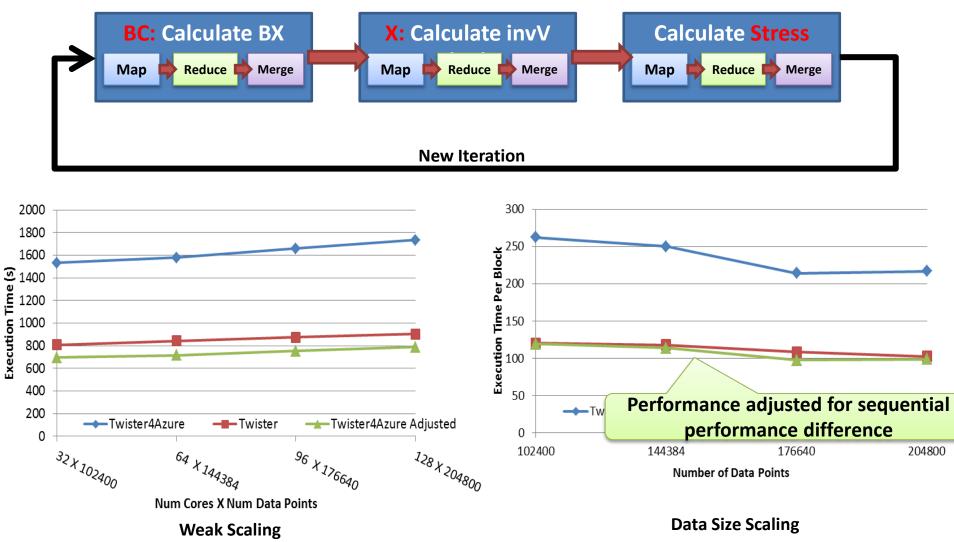
I. Borg and P. J. Groenen. *Modern Multidimensional Scaling: Theory and Applications. Springer, New* York, NY, U.S.A., 2005.
C. Bishop, M. Svens'en, and C. Williams. GTM: The generative topographic mapping. *Neural computation*, 10(1):215–234, 1998.

# **Multidimensional Scaling**

- Scaling by Majorizing a Complicated Function
- Can be merged to Kmeans result



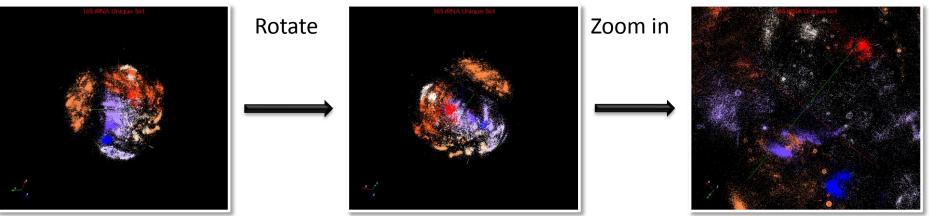
### Multi Dimensional Scaling on Twister (Linux), Twister4Azure and Hadoop



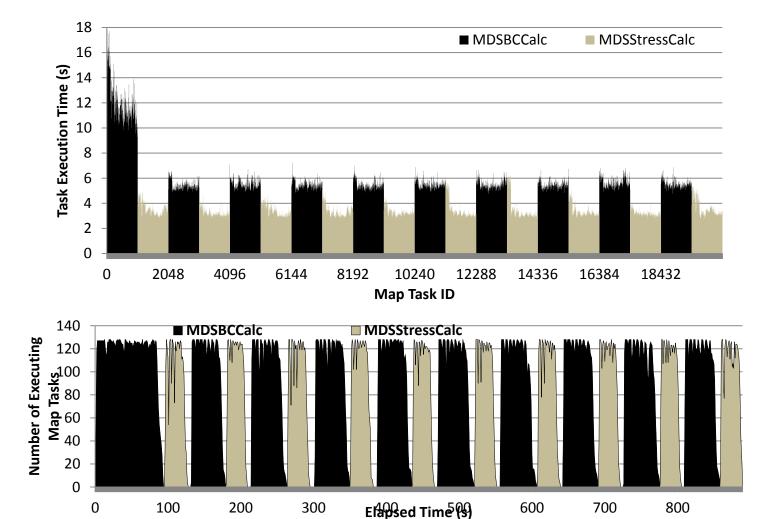
Scalable Parallel Scientific Computing Using Twister4Azure. Thilina Gunarathne, BingJing Zang, Tak-Lon Wu and Judy Qiu. Submitted to Journal of Future Generation Computer Systems. (Invited as one of the best 6 papers of UCC 2011)

# Visualization

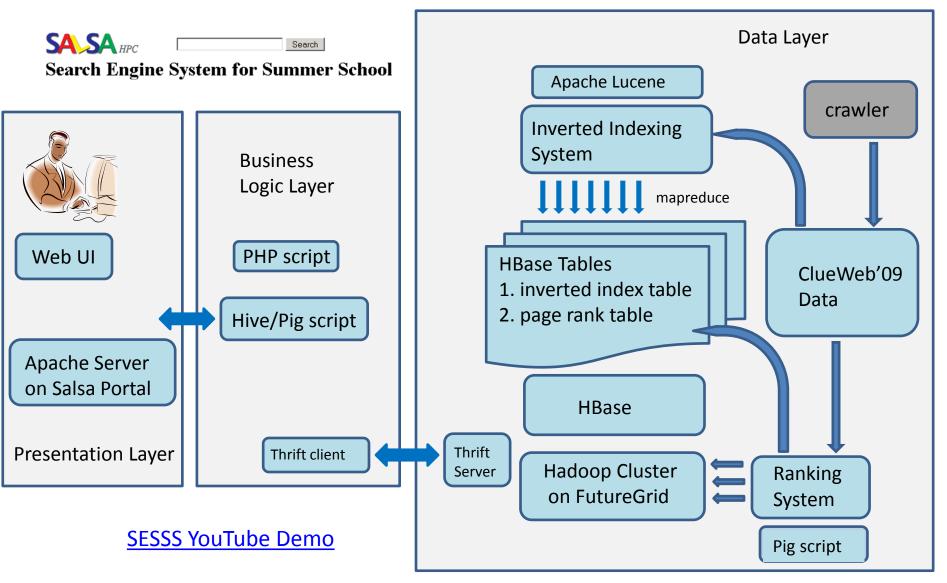
- Used PlotViz3 to visualize the 3D plot generated in this project
- It can show the sequence name, highlight interesting points, even remotely connect to HPC cluster and do dimension reduction and streaming back result.



#### Multi Dimensional Scaling on Azure



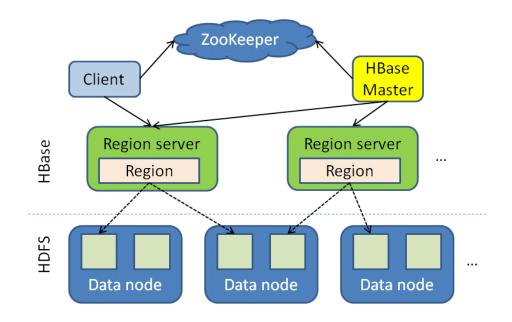
## **Architecture for Search Engine**



## Parallel Inverted Index using HBase

- Get inverted index involved in HBase "cloud" -> doc1, doc2, ... "computing" -> doc1, doc3, ...
- 1. Store inverted indices in HBase tables scalability and availability
- 2. Parallel index building with MapReduce (supporting Twister doing data mining on top of this)
- 3. Real-time document insertion and indexing
- 4. Parallel data analysis over text as well as index data
- 5. ClueWeb09 data set for experiments in an HPC environment

### HBase architecture:



- Tables split into regions and served by region servers
- Reliable data storage and efficient access to TBs or PBs of data, successful application in Facebook and Twitter
- Problem: no inherent mechanism for field value searching, especially for full-text values

## ClueWeb09 dataset

- Whole dataset: about 1 billion web pages in ten languages collected in 2009
- Category B subset:

# of web pages	Language	# of unique URLs	Compressed size	Uncompressed size
50 million	English	4,780,950,903	250GB	1.5TB

- Data stored in .warc.gz files, file size : 30MB 200MB
- Major fields in a WARC record:
  - HTML header record type, e.g., "response"

- TREC ID: a unique ID in the whole dataset, e.g., "clueweb09en0040-54-00000"

- Target URL: URL of the web page
- Content: HTML page content

## Table schemas in HBase

#### • Data table schema for storing the ClueWeb09 data set:



• Index table schema for storing term frequencies:

	frequencies			
	"283"	"1349"	(other document ids)	
"database" —>	3	4		

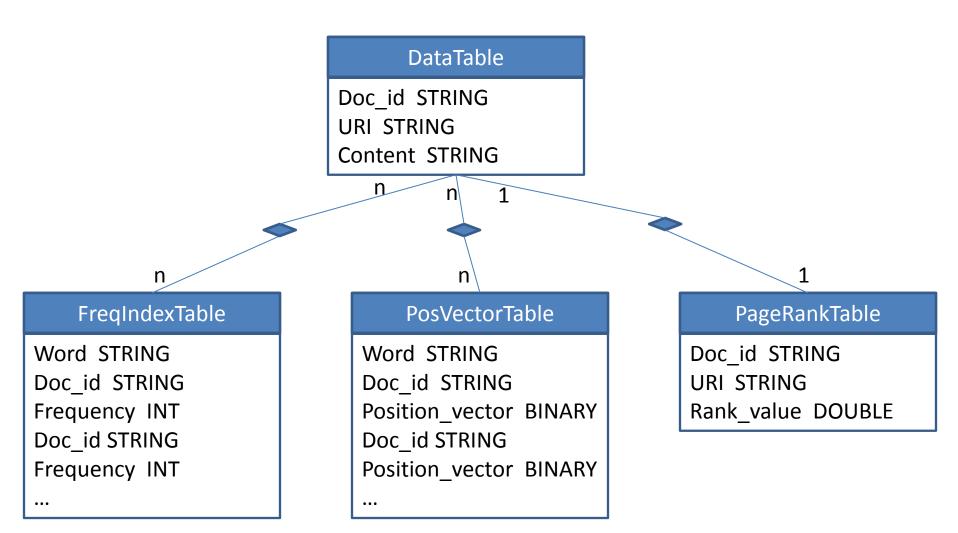
• Index table schema for storing term position vectors:

	positions		
	"283"	"1349"	(other document ids)
"database" →	1, 24, 33	1, 34, 77, 221	

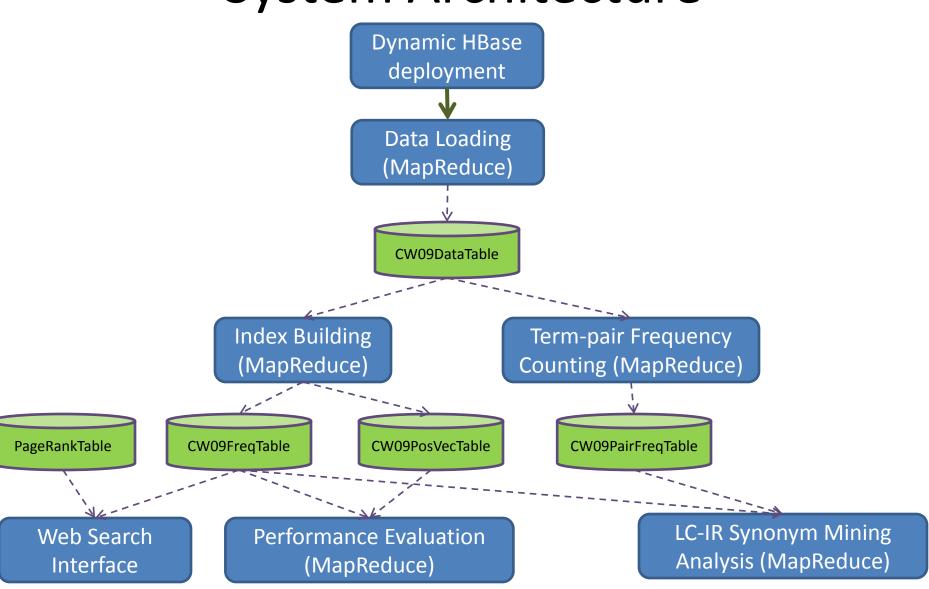
• Table schema for PageRank values:

	PageRanks		
	URI	RankValue	
"20000001">	http://en.wikipedia.org/wiki/	43.6	

### Table schemas – Entity Relation Diagram



## System Architecture



# LC-IR Synonym Mining

 Mining synonyms from large document sets based on words' co-appearances

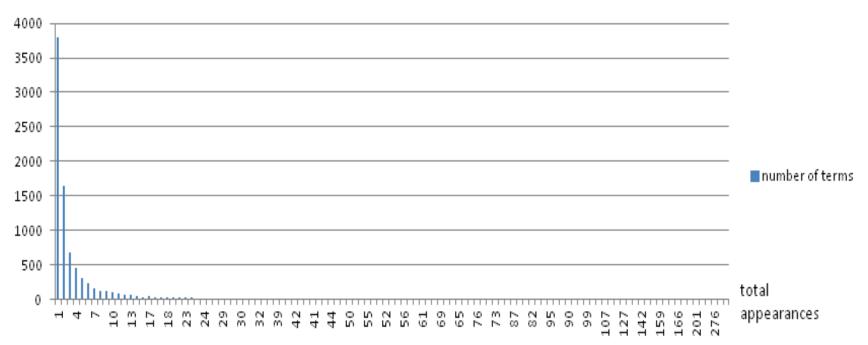
 $Similarity_{\text{LC}-\text{IR}}(w_1, w_2) = \frac{\min(\text{Hits}(``w_1 w_2"), \text{Hits}(``w_2 w_1"))}{\text{Hits}(w_1) \times \text{Hits}(w_2)}$ 

Steps for completing LC-IR synonym mining in HBase:
1. Scan the data table and generate a "pair count" table for word-pairs;

2. Scan the "pair count" table and calculate similarities, looking up single word hits in the index table;

3. Filter the pairs with similarities lower than a threshold.

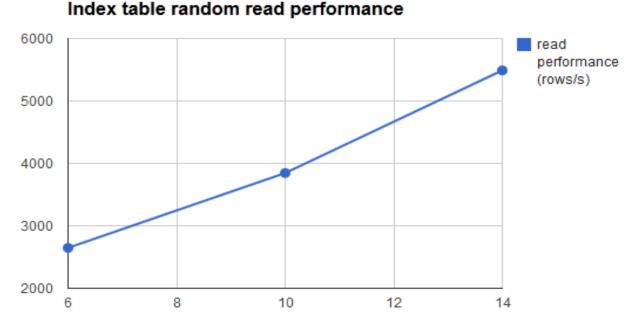
Distribution of total appearances in all documents



- 100 documents indexed, 8499 unique terms
- 3793 (45%) terms appear only once in all documents
- Most frequent word: "you"

- Example synonyms mined (among 16516 documents):
  - chiropodists podiatrists (0.125, doctors for foot disease)
  - desflurane isoflurane (0.111, narcotic)
  - dynein kinesin (0.111, same type of protein)
  - menba monpa (0.125, a nation/race of Chinese people living in Tibet)
  - lyrica pregabalin (0.125, different names for the same medicine for diabetes)

- Preliminary performance evaluation
  - 6 distributed clients started, each reading 60000 random rows
  - average speed: 2647 rows/s



number of distributed clients

- Original data table size: 29GB (2,594,536 documents)
- Index table size: 8,557,702 rows (one row for each indexed term)
- Largest row: 2,580,938 cell values, 162MB uncompressed size
- At most 1000 cell values are read from each row in this test
- Aggregate read performance increases as number of concurrent clients increases

Number of nodes	Number of mappers	Index building time (seconds)
8	32	18590
12	37 (15.6% increase)	16142 (15.2% improvement)
16	47 (46.9% increase)	13480 (37.9% improvement)

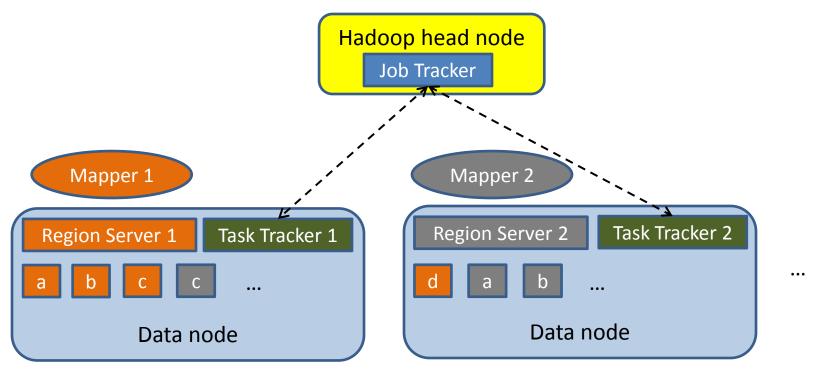
Index building performance vs. resources increase

- Original data table size: 29GB (2,594,536 documents)
- 6 computing slots on each node
- HBase overhead: data transmission to region servers, cell value sorting based on keys, gzip compression/decompression
- Number of mappers not doubled when number of nodes doubled because of small table size
- Increase in index building performance is close to increase in number of mappers

## **Practical Problems and experiences**

- Hadoop and HBase configuration
  - Lack of "append" support in some versions of Hadoop: missing data, various errors in HBase and HDFS.
  - Low data locality in HBase MapReduce: "c046.cm.cluster" for Task Tracker vs. "c046.cm.cluster." for Region Server.
  - Clock not synchronized error: clock not synched with NTP on some nodes.
- Optimizations in the synonym mining programs
  - Addition of a word count table with bloom filter.
  - Local combiners for word pair counter.
  - Caching of word counts during the synonym scoring phase.

#### Low data locality in MapReduce over HBase



- Data splits assigned to mappers by regions (one mapper per region in most cases)
- Mapper deployment based on mapper-region server locality
- Problem: region data blocks not necessarily local to region servers
- Data locality gets even worse after region splits or region server failures

# Ackowledgements

#### SALSA HPC Group

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